

**ALTERA.**  
UNIVERSITY  
PROGRAM

**5G**

**The 1st 5G Algorithm  
Innovation Competition**

**SCMA**

Organizer |



西安电子科技大学  
XIDIAN UNIVERSITY

**terasic** Terasic Inc.

Sponsor |



**SPREADTRUM**

# 1st 5G Algorithm Innovation Competition-SCMA

Task	Description	Requirements
SCMA- sparse code multiple access for the next generation wireless communic ations	<p>Multiple access is among the core physical layer technologies of wireless communications, which enables wireless base stations to identify a large number of different terminal users and serve them simultaneously. Current systems choose to use orthogonal multiple access method, i.e., users are orthogonal to each other in at least one radio resource dimension (e.g., frequency, time, code and etc.) . The OFDMA technology used in 4G systems is one example, in which radio resources is divided into two-dimensional time-frequency grids and each grid can only be used by one user at a time.</p> <p>It is obvious that the number of simultaneously accessible user is strictly proportional to the number of available orthogonal resources, and is thus limited. Facing the 5G requirement of massive connectivity, non-orthogonal multiple access becomes the research focus of 5G multiple access technologies, among which, Sparse code multiple access (SCMA) is a promising candidate. At the transmitter, coded bits are directly mapped to multi-dimensional codewords in complex domain, and codewords from different users are overlapped non-orthogonally in a sparse spreading way; the receiver performs joint multiuser detection followed by channel decoding for data recovering. Thanks to the sparsity, low complexity algorithms could be design to achieve near optimal detection results. The number of non-orthogonally superposed codewords can be much larger than the number of orthogonal resource units. This leads to the advantage of SCMA to serve more users while keeping the same expanse of resources, thus effectively improve the overall system capacity.</p> <p>According to the introduction of SCMA encoding and decoding principles in the given material, please design and then implement a simplified SCMA uplink communication system, and verify the feasibility and performance of such non-orthogonal multiple access system by software simulations and hardware testing.</p>	<p><b>Requirements:</b> According to the introduction given in the SCMA training materials, please design and then implement a simplified SCMA uplink multiple access communication system (abstract system, rather than a complete communication system) , focusing especially on the development of SCMA encoder and the corresponding low complexity decoder design.</p> <p><b>Works format:</b></p> <ol style="list-style-type: none"><li>1. Complete the detailed system design documents for the simplified SCMA uplink multiple access system, especially the low-complexity SCMA decoder design.</li><li>2. Complete Matlab simulation for the simplified SCMA uplink multiple access system, and gives BER vs Eb / No performance curve.</li><li>3. Complete FPGA logic design and implementation of the SCMA uplink multiple access system, test its performance, compare with simulation curves and report the resource usage.</li></ol> <p><b>Delivery Material:</b></p> <ol style="list-style-type: none"><li>1. SCMA uplink multi-access system design documents, code and simulation results;</li><li>2. FPGA design specifications, code, and test results of the bit file.</li></ol> <p><b>Selection criteria in the first round:</b></p> <ol style="list-style-type: none"><li>1. Correctly understand SCMA system, complete the detailed system design documents for the simplified SCMA uplink multiple access system, especially the low-complexity decoder.</li><li>2. Complete the link-level Matlab / C simulations for the simplified system, and give BER vs Eb / No performance curve .</li><li>3. Complete the framework and methodology design of using the given FPGA platform for the SCMA system implementation.</li></ol> <p><b>Selection criteria in the second round:</b></p> <ol style="list-style-type: none"><li>1. Complete the FPGA logic design document, and implement the simplify SCMA uplink multiple access system with given FPGA platform, achieving a correct decoding rate more than 99.9%.</li><li>2. Test the BER vs Eb / No performance curve, and compare the results with simulation (allow refreshed optimization in simulations, the difference should be less than 1dB</li><li>3. FPGA implementation of the low complexity SCMA decoder, data throughput, processing delays and FPGA chip resource should be optimized.</li></ol>

# References to Read

## \_Helpful Documents for Understanding SCMA

### MUST Read Papers:

1. 《SCMA Codebook Design》 (to understand SCMA)  
<http://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=6966170>
2. 《Novel low-density signature for synchronous cdma systems over AWGN channel》 (to understand MPA)  
<http://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=4471881>

### SCMA Related Publications:

- H. Nikopour and H. Baligh, “Sparse Code Multiple Access,” IEEE PIMRC, 2013.
- M. Taherzadeh, H. Nikopour, A. Bayesteh, and H. Baligh, “SCMA Codebook Design,” IEEE VTC-fall, 2014.
- Kelvin Au, Liqing Zhang, Hosein Nikopour, Eric Yi, Alireza Bayesteh, Usa Vilaipornsawai, Jianglei Ma, Peiying Zhu, “Uplink Contention Based SCMA for 5G Radio Access,” IEEE Globecom 5G workshop 2014.
- H. Nikopour, E. Yi, A. Bayesteh, K. Au, M. Hawryluck, H. Baligh, and Jianglei Ma, “SCMA for Downlink Multiple Access of 5G Wireless Networks,” IEEE Globecom 2014.
- S. Zhang, X. Xu, L. Lu, Y. Wu, G. He, and Y. Chen, “Sparse Code Multiple Access: An Energy Efficient Uplink Approach for 5G Wireless Systems,” IEEE Globecom 2014.
- A. Bayesteh, E. Yi, E. , H. Nikopour, H. Baligh, “Blind Detection of SCMA for Uplink Grant-Free Multiple-Access”, ISWCS 2014.
- Y. Wu, S. Zhang, and Y. Chen, “Iterative multiuser receiver in sparse code multiple access systems,” IEEE ICC 2015.

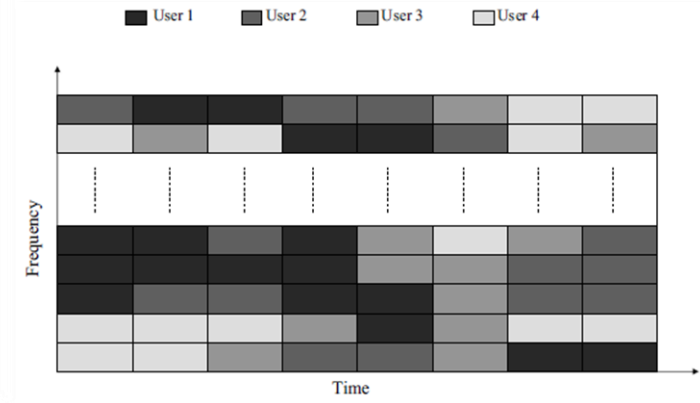
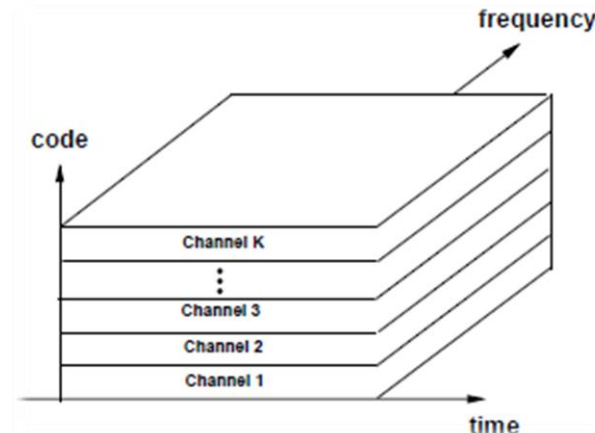
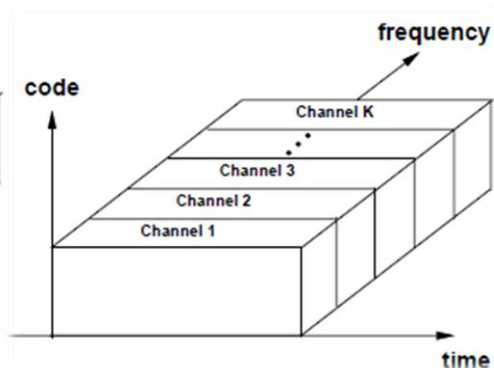
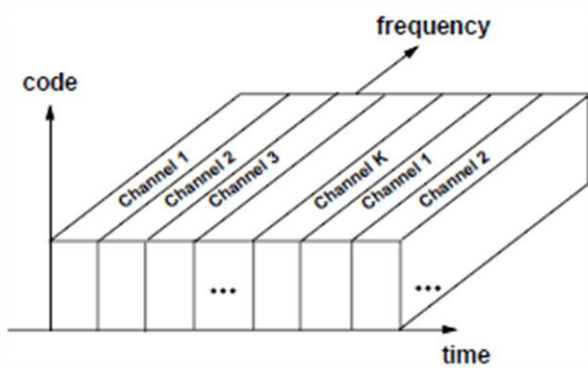
# Outline

- What is SCMA?
- Why we need SCMA in 5G?
- How does SCMA work?
- What will you implement?

To have a gut feeling what is sparse code multiple access

## **WHAT IS SCMA?**

# Existing Multiple Access Schemes



## TDMA/FDMA

- 2G Communication system, e.g. GSM
- Orthogonal in time or frequency domain
- Users are scheduled on orthogonal time slots

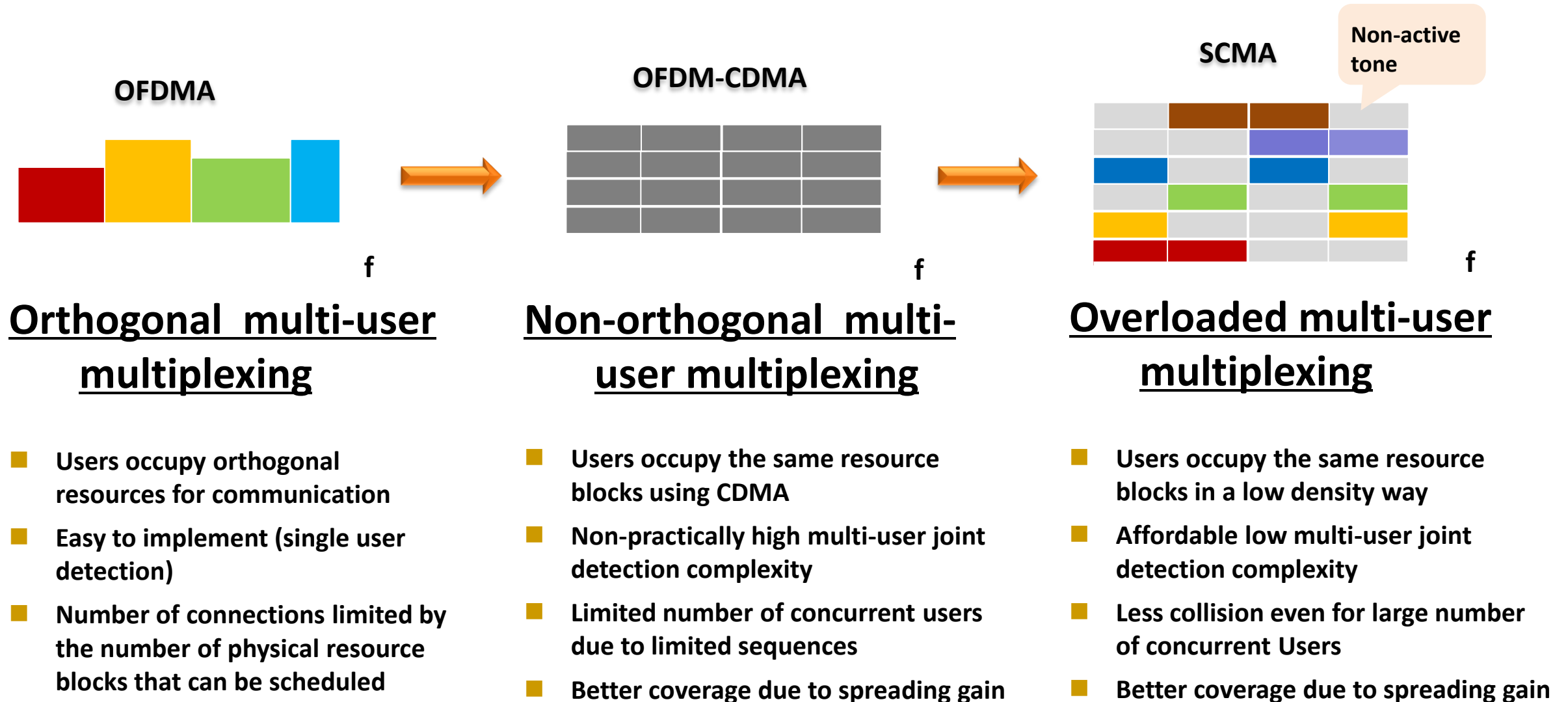
## CDMA

- 3G Communication system, e.g. WCDMA
- Non-orthogonal in time and frequency but orthogonal in code domain
- Users are scheduled on orthogonal sequences

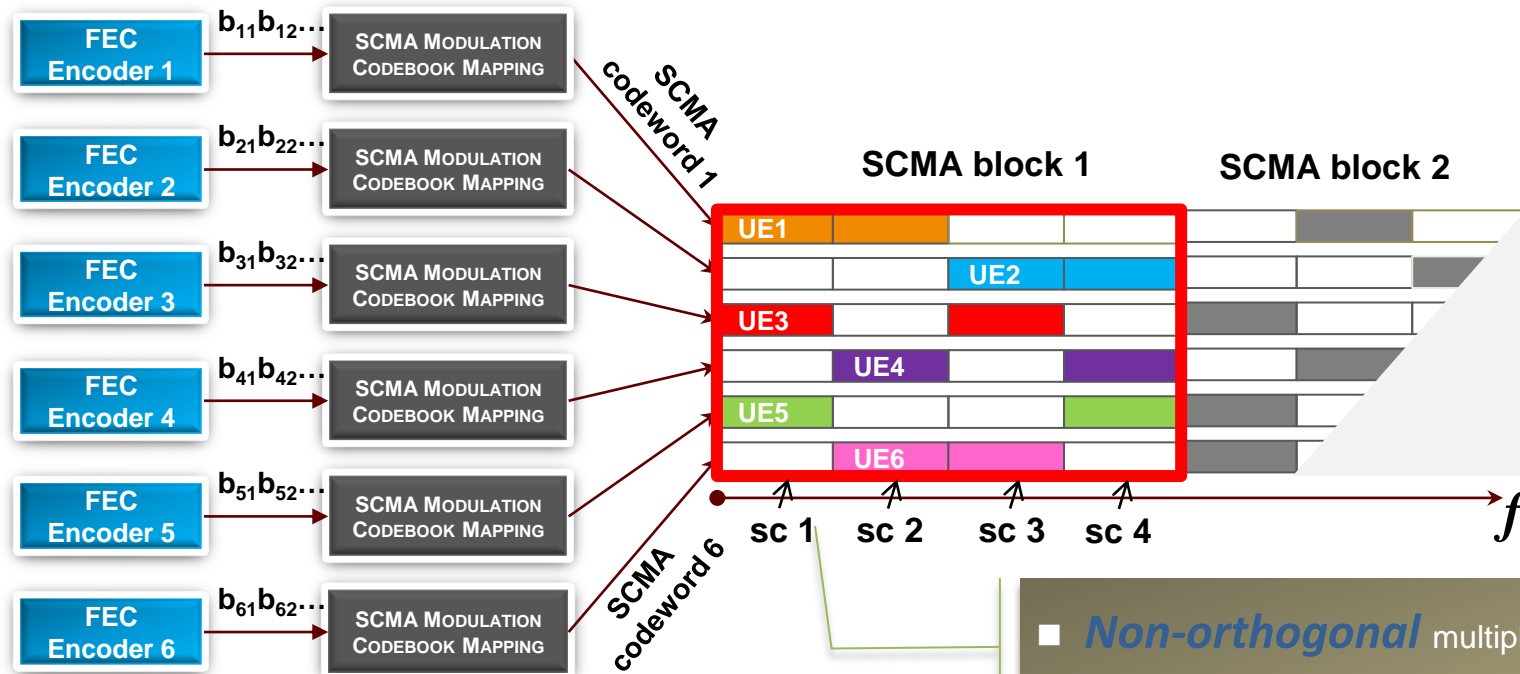
## OFDMA

- 4G Communication system, e.g. LTE
- Orthogonal in 2D time-frequency lattice domain
- Users are scheduled on orthogonal time-frequency lattice

# From OFDMA to SCMA



# SCMA (Sparse Code Multiple Access)



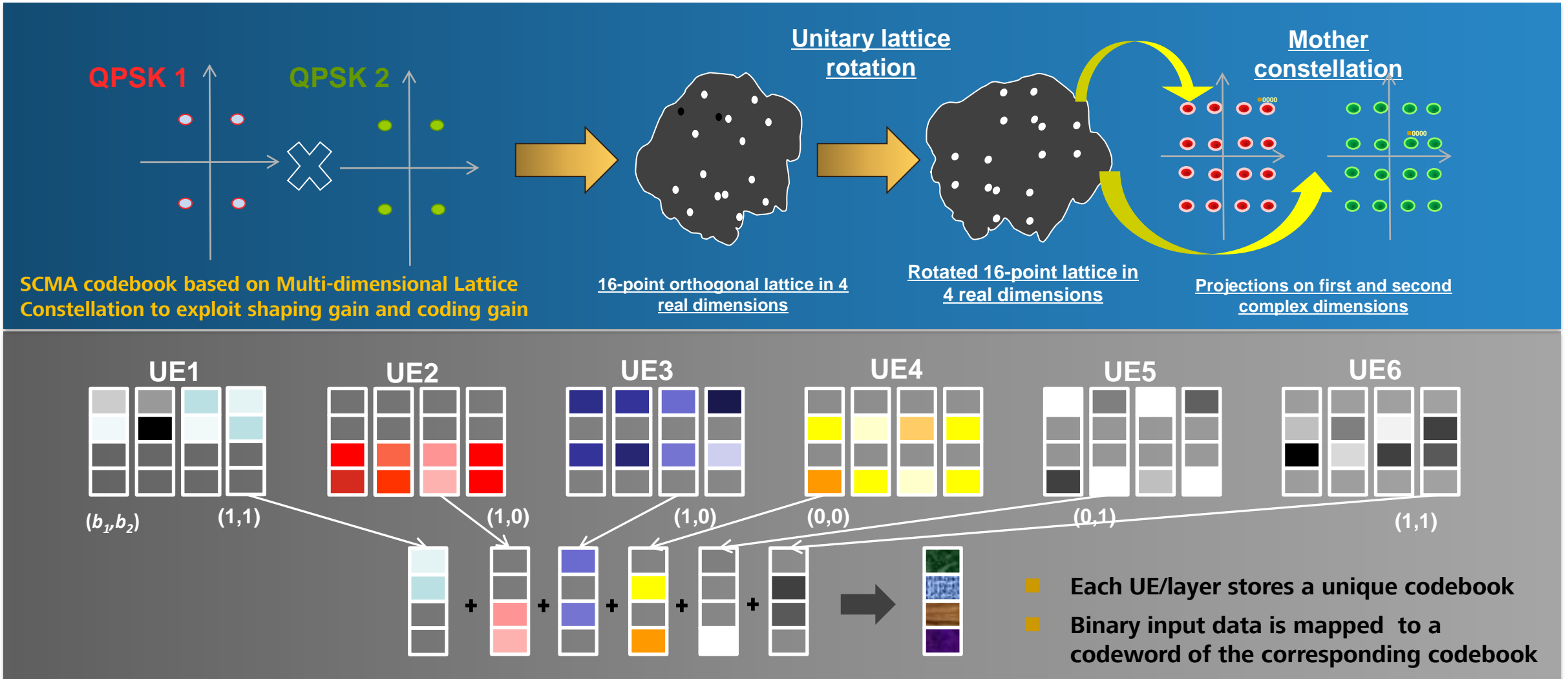
## A new frequency domain non-orthogonal waveform

- ❑ Input bits are directly mapped to codewords and spread over multiple sub-carriers
- ❑ Codewords can be assigned to same UE or different UEs

- **Non-orthogonal** multiplexing of code layers
- **Over-Loading** to increase overall rate and connectivity
- **Sparsity** to limit Rx complexity for detection
- **Spreading** for robust link-adaptation, coverage
- **Multi-dimensional** codewords with shaping gain



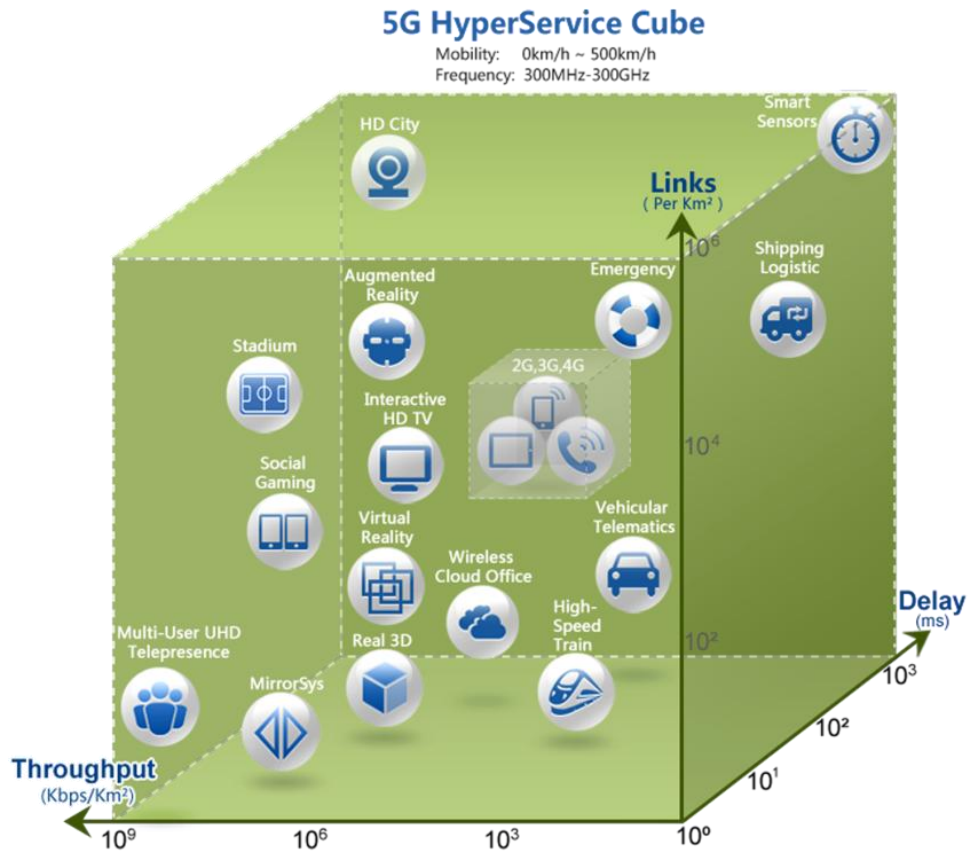
# SCMA Codebook Design



To know what role SCMA plays in 5G and what benefit it brings along

## **WHY WE NEED SCMA IN 5G?**

# 5G Vision: Zero Distance Communications



**Massive Capacity**



**Massive Connectivity**



**Zero Waiting**

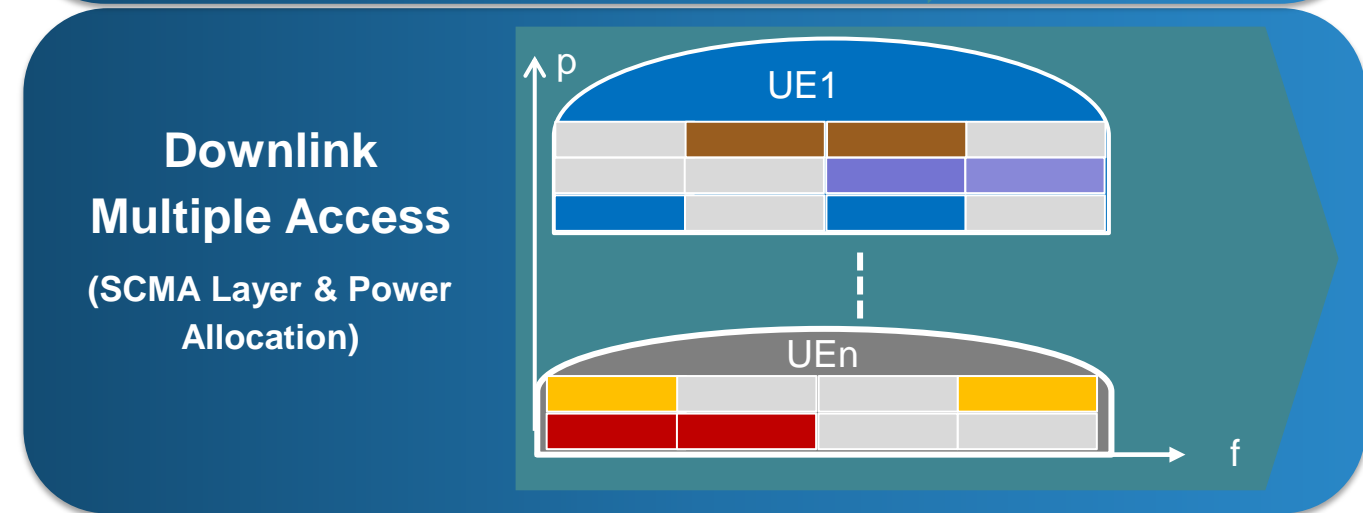
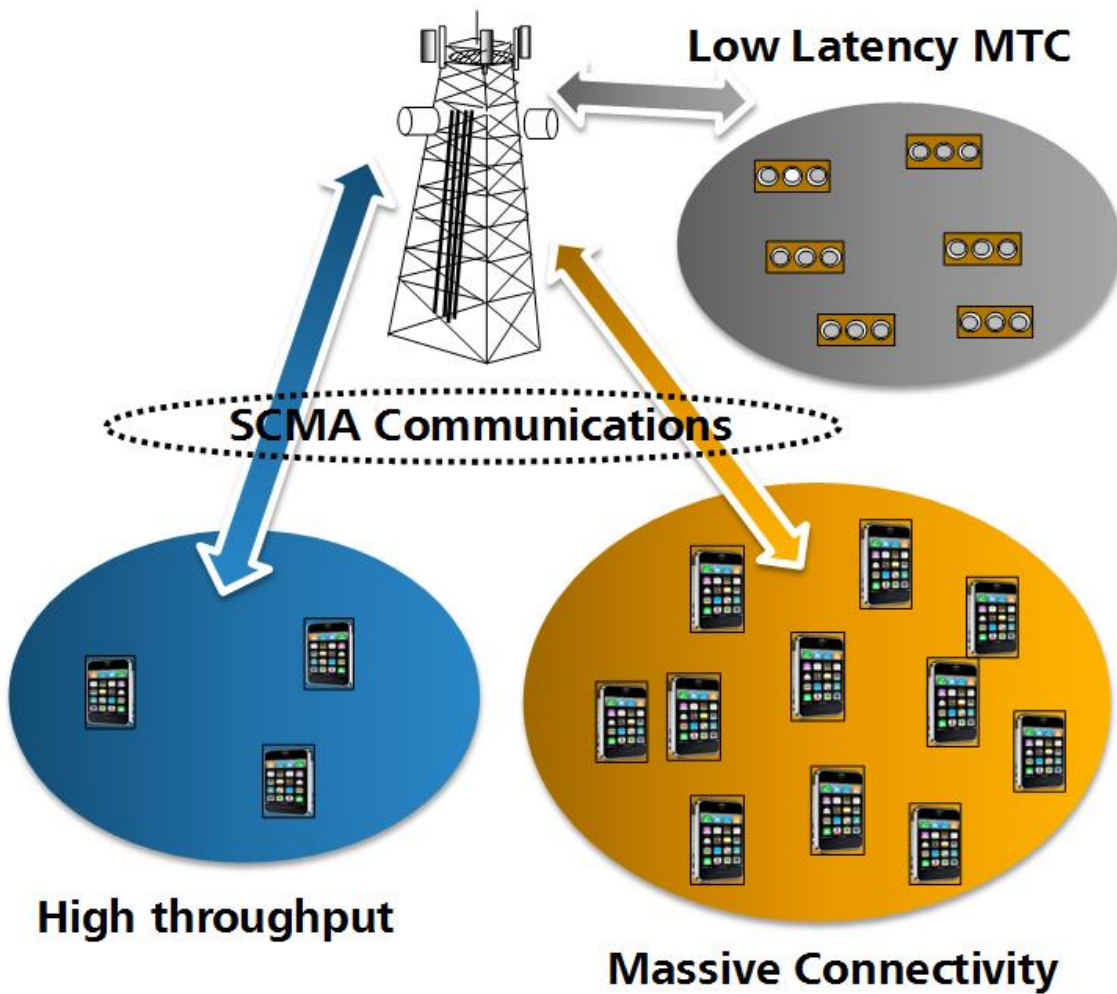


**10<sup>3</sup>**  
 Times Traffic Flooding

**10<sup>2</sup>**  
 Billion Connections

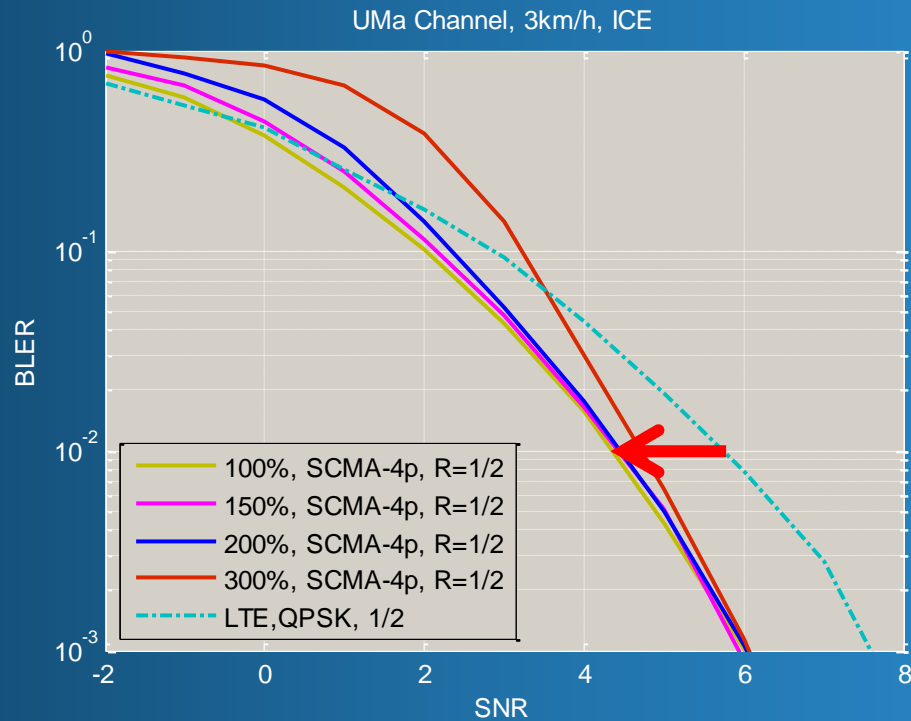
**10<sup>0</sup>**  
 ms Delay Experience

# Example of SCMA Application Scenarios



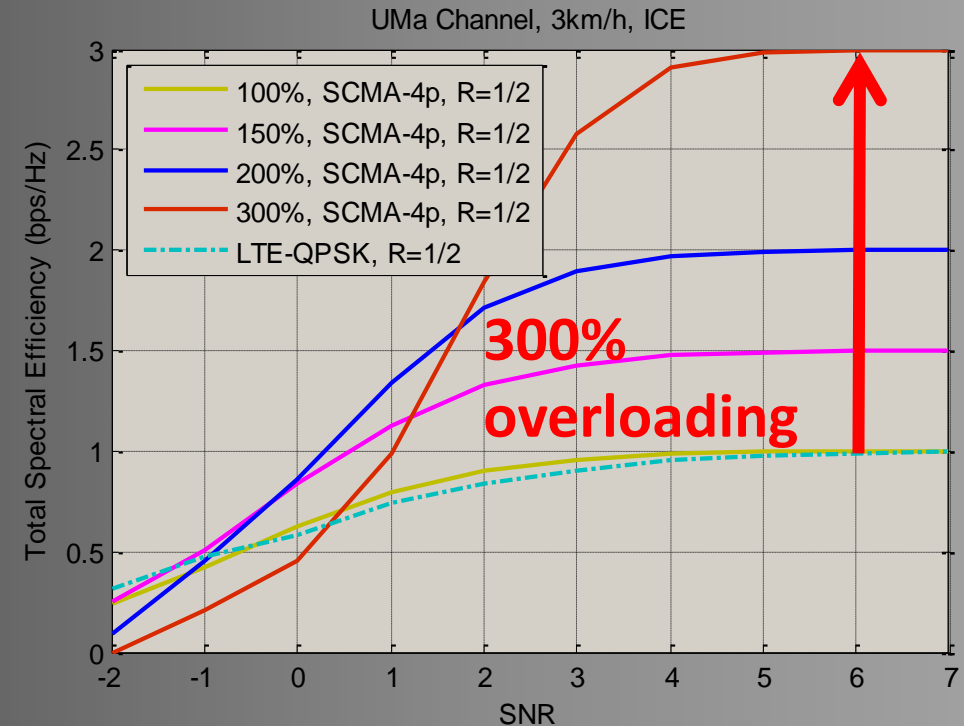
# SCMA Offers Better Link Quality and 300% Larger Number of Physical Link Connections over LTE

## Better Link Quality over LTE



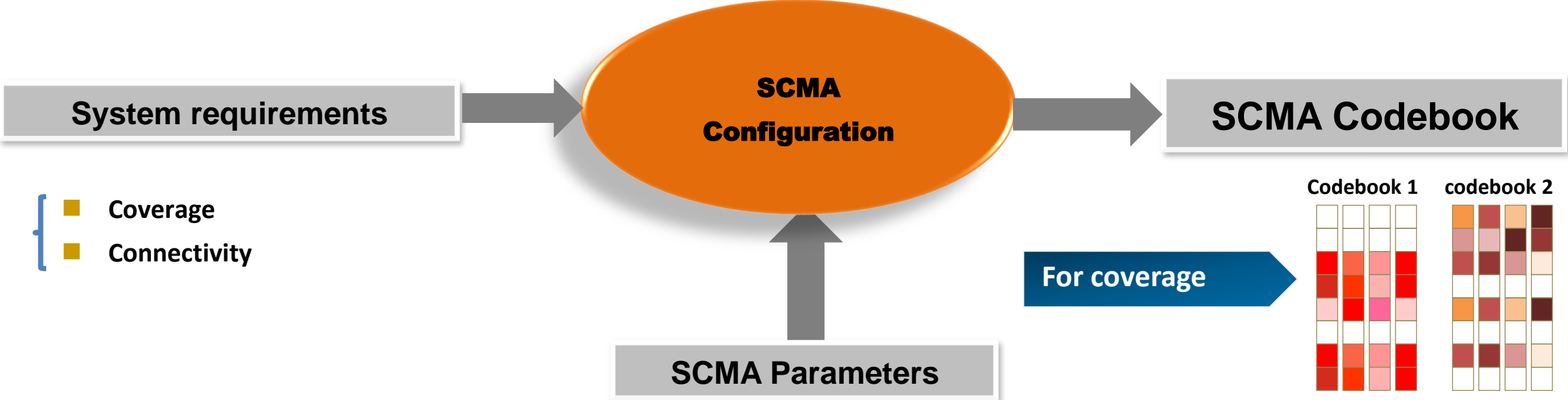
- SCMA has SNR gain over LTE (same rate & power per user)
- SCMA with overloading performs towards single user

## 300% Larger Number of Connected Users



- Given the same SNR, SCMA can boost total system throughput up to 300% over LTE (@LTE BLER=0.01)

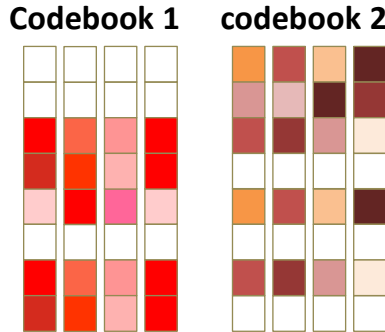
# SCMA Codebook Design Can Flexibly Adapt to Meet Diversified System Requirements



- Coverage
- Connectivity

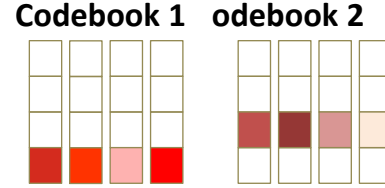
- Number of codewords of an SCMA codebook:  $M$
- Spreading factor:  $F$
- Max number of layers (or codebooks/signatures) :  $V$
- Number of nonzero elements of each codeword:  $J$

For coverage



Example: larger SF, more non-zero elements

For connectivity



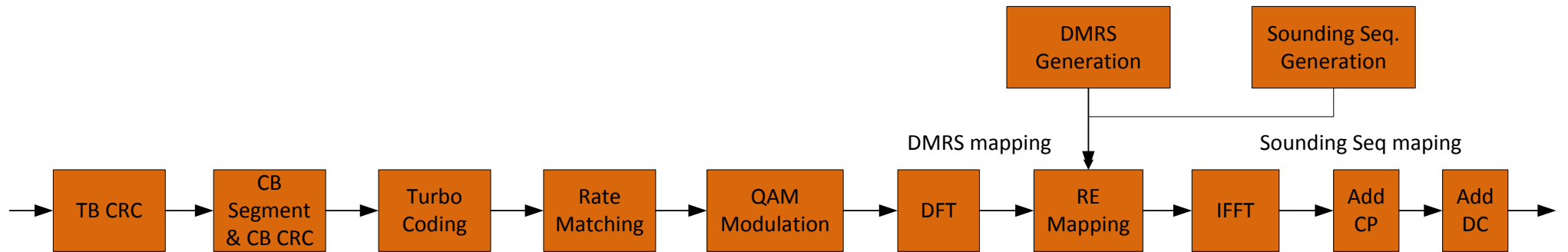
Example: moderate SF, one non-zero element

To have a gut feeling how SCMA will be implemented in the 5G wireless systems

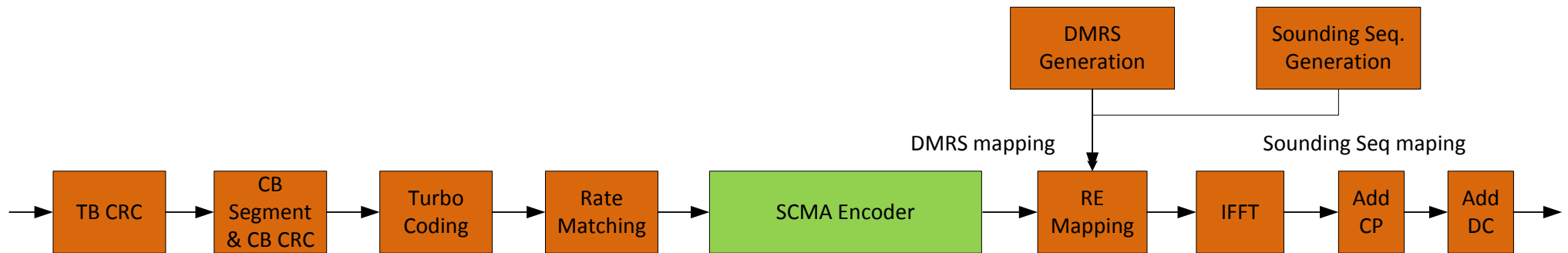
## **HOW DOES SCMA WORK?**

# SCMA Uplink Transmission System Diagram

LTE SC-FDMA Transmitter



SCMA over OFDM Transmitter

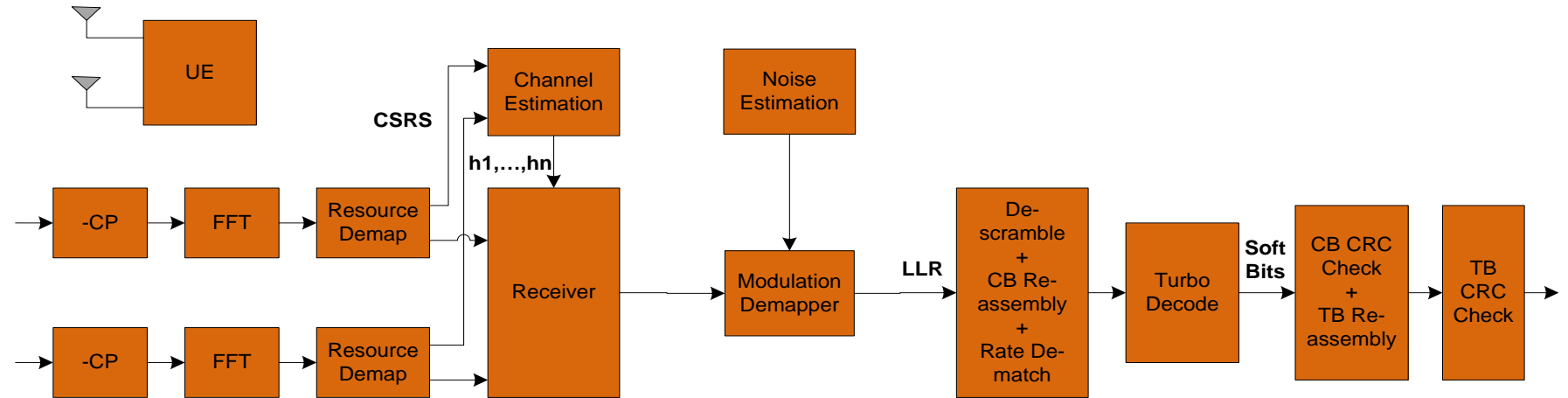


Wanna know more about 4G LTE? Please refer to 3GPP standards or the book "LTE the UMTS Long Term Evolution, from Theory to Practice" from Wiley Press.

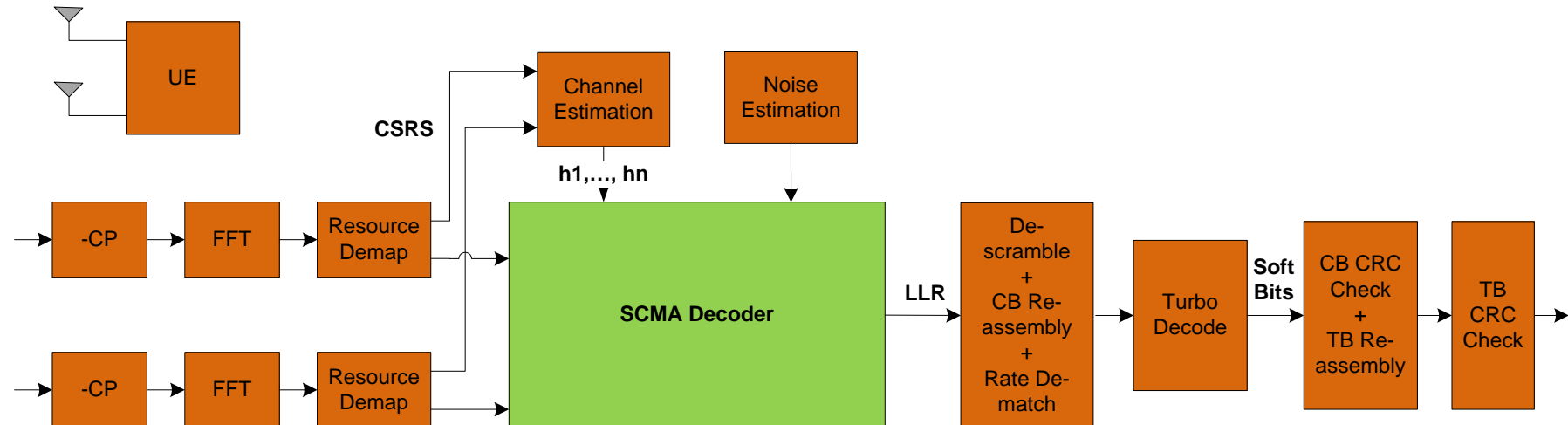


# SCMA Uplink Transmission System Diagram

LTE OFDMA Receiver



SCMA over OFDM Receiver



# Example of SCMA Codebook

## Codebook Related Parameters

Related Variables	Typical value	Description
V	6	6 variable nodes (VN), number of data layers
F	4	4 function nodes (FN), number of physical resources
d <sub>f</sub>	3	Each FN is connected to 3 VNs
d <sub>v</sub>	2	Each VN is connected to 2 FNs
M	4	Number of codeword in each codebook
CB <sub>i</sub>	F-by-M matrix	Codebook for one SCMA data layer

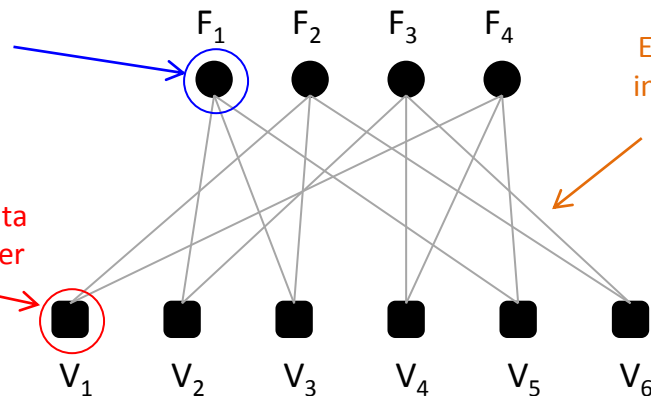
## Codebook in Storage (V=6, F=4, df=3, dv=2, M=4)

SCMA Codebook index	SCMA codebook for each layer
CB <sub>1</sub>	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ -0.1815 - 0.1318i & -0.6351 - 0.4615i & 0.6351 + 0.4615i & 0.1815 + 0.1318i \\ 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \end{bmatrix}$
CB <sub>2</sub>	$\begin{bmatrix} 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ 0 & 0 & 0 & 0 \\ -0.1815 - 0.1318i & -0.6351 - 0.4615i & 0.6351 + 0.4615i & 0.1815 + 0.1318i \\ 0 & 0 & 0 & 0 \end{bmatrix}$
CB <sub>3</sub>	$\begin{bmatrix} -0.6351 + 0.4615i & 0.1815 - 0.1318i & -0.1815 + 0.1318i & 0.6351 - 0.4615i \\ 0.1392 - 0.1759i & 0.4873 - 0.6156i & -0.4873 + 0.6156i & -0.1392 + 0.1759i \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$
CB <sub>4</sub>	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ -0.0055 - 0.2242i & -0.0193 - 0.7848i & 0.0193 + 0.7848i & 0.0055 + 0.2242i \end{bmatrix}$
CB <sub>5</sub>	$\begin{bmatrix} -0.0055 - 0.2242i & -0.0193 - 0.7848i & 0.0193 + 0.7848i & 0.0055 + 0.2242i \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ -0.6351 + 0.4615i & 0.1815 - 0.1318i & -0.1815 + 0.1318i & 0.6351 - 0.4615i \end{bmatrix}$
CB <sub>6</sub>	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ 0.1392 - 0.1759i & 0.4873 - 0.6156i & -0.4873 + 0.6156i & -0.1392 + 0.1759i \\ 0 & 0 & 0 & 0 \end{bmatrix}$

## Tanner Graph Representation

Function node, representing the physical resource elements (PREs)

Variable node, representing the data from one SCMA layer



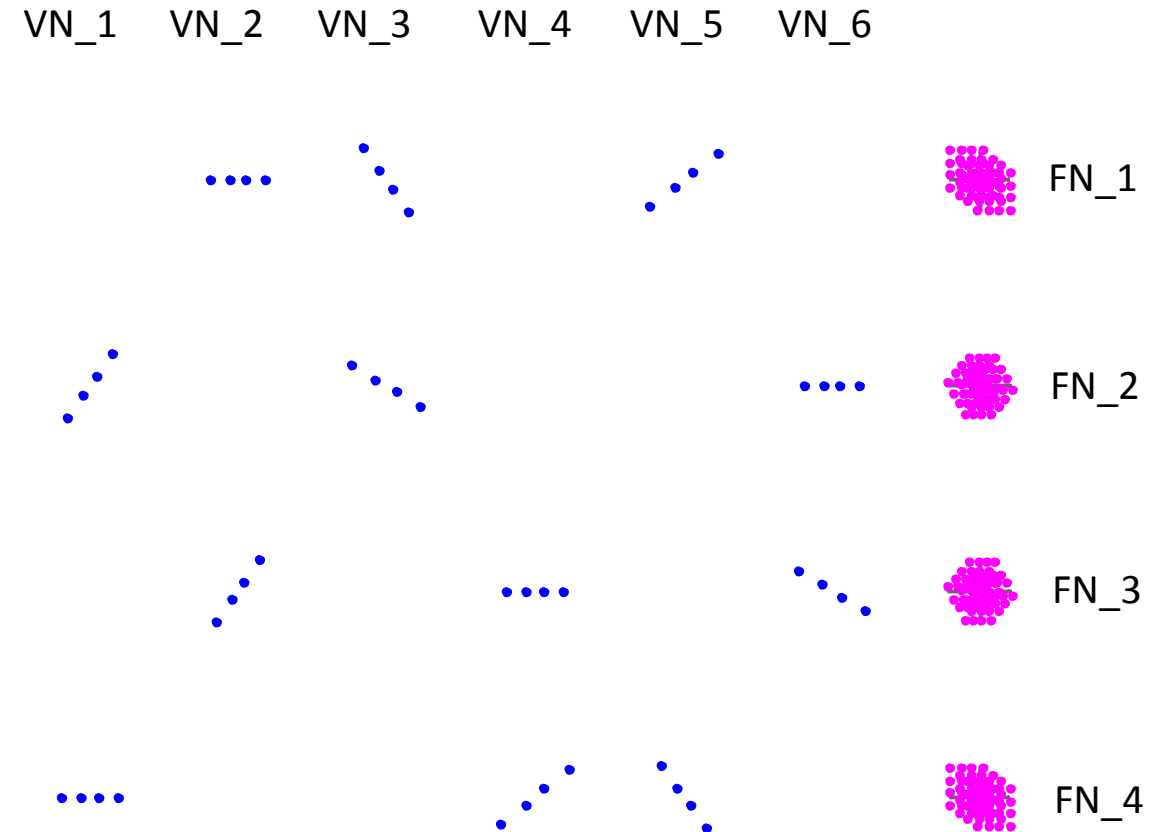
Edge for passing the inference of the data symbols

# Example of SCMA Codebook

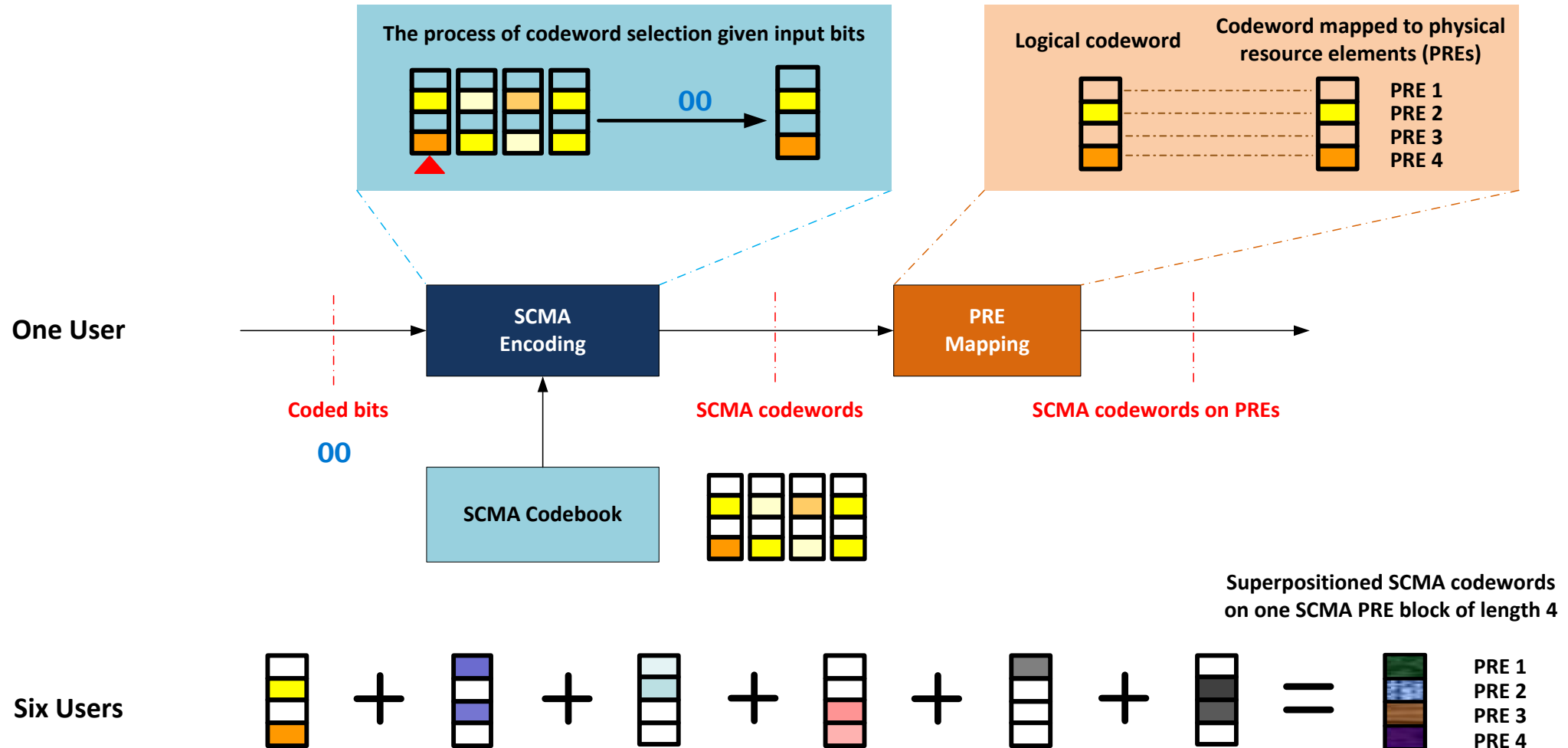
Codebook in Storage (V=6, F=4, df=3, dv=2, M=4)

SCMA Codebook index	SCMA codebook for each layer
CB_1	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ -0.1815 - 0.1318i & -0.6351 - 0.4615i & 0.6351 + 0.4615i & 0.1815 + 0.1318i \\ 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \end{bmatrix}$
CB_2	$\begin{bmatrix} 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ 0 & 0 & 0 & 0 \\ -0.1815 - 0.1318i & -0.6351 - 0.4615i & 0.6351 + 0.4615i & 0.1815 + 0.1318i \\ 0 & 0 & 0 & 0 \end{bmatrix}$
CB_3	$\begin{bmatrix} -0.6351 + 0.4615i & 0.1815 - 0.1318i & -0.1815 + 0.1318i & 0.6351 - 0.4615i \\ 0.1392 - 0.1759i & 0.4873 - 0.6156i & -0.4873 + 0.6156i & -0.1392 + 0.1759i \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$
CB_4	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ -0.0055 - 0.2242i & -0.0193 - 0.7848i & 0.0193 + 0.7848i & 0.0055 + 0.2242i \end{bmatrix}$
CB_5	$\begin{bmatrix} -0.0055 - 0.2242i & -0.0193 - 0.7848i & 0.0193 + 0.7848i & 0.0055 + 0.2242i \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ -0.6351 + 0.4615i & 0.1815 - 0.1318i & -0.1815 + 0.1318i & 0.6351 - 0.4615i \end{bmatrix}$
CB_6	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0.7851 & -0.2243 & 0.2243 & -0.7851 \\ 0.1392 - 0.1759i & 0.4873 - 0.6156i & -0.4873 + 0.6156i & -0.1392 + 0.1759i \\ 0 & 0 & 0 & 0 \end{bmatrix}$

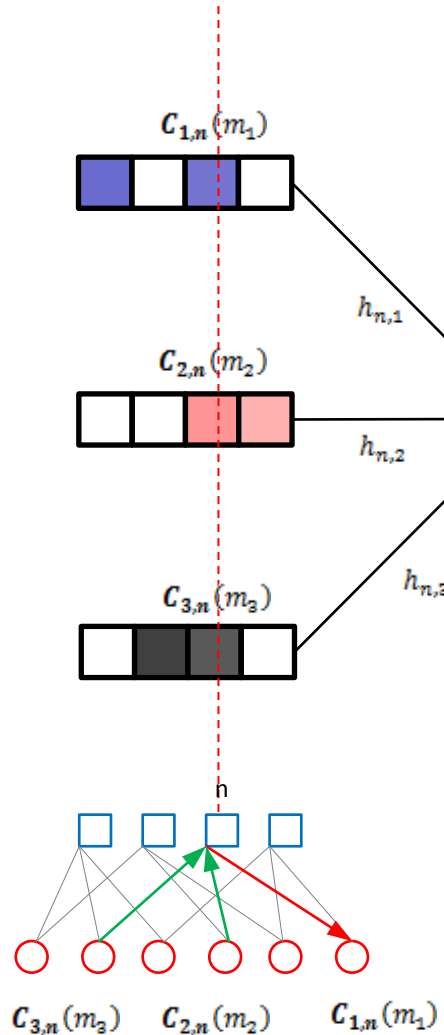
Codebook Presented by Constellation Points



# How to Do SCMA Encoding with SCMA Codebook



# How to Do SCMA Decoding



$C_{k,n}(m_k)$ $n=1, \dots, F$ $k=1, \dots, V$	The constellation symbol of VN node $k$ on physical resource $n$ when using codeword $m_k$
$h_{n,k}$ $n=1, \dots, F$ $k=1, \dots, V$	Channel coefficient of user $k$ (layer $k$ / VN node $k$ ) on physical resource $n$
$z_n$ $n=1, \dots, F$	White noise at physical resource $n$
$y_n$ , $n=1, \dots, F$	Received signal as input to the MPA decoder on resource $n$

- Optimal **Maximum joint A posteriori Probability (MAP)** detection

$$\hat{\mathbf{x}} = \arg \max_{\mathbf{x} \in \mathcal{X}^K} p(\mathbf{x}|\mathbf{y}).$$

$$\hat{x}_k = \arg \max_{a \in \mathcal{X}} \sum_{\substack{\mathbf{x} \in \mathcal{X}^K \\ x_k = a}} p(\mathbf{x}|\mathbf{y}), \quad \forall k$$

- Equivalence of MAP with **Maximum Likelihood (ML)** detection when the a prior probabilities of  $x_k$  are the same

$$p(\mathbf{x}|\mathbf{y}) = \frac{p(\mathbf{y}|\mathbf{x})P(\mathbf{x})}{P(\mathbf{y})} \propto p(\mathbf{y}|\mathbf{x})P(\mathbf{x})$$

$$P(\mathbf{y}) = \sum_{\mathbf{x} \in \mathcal{X}^K} p(\mathbf{y}|\mathbf{x})P(\mathbf{x})$$

$$P(\mathbf{x}) = \prod_{k=1}^K P(x_k)$$

$$p(\mathbf{y}|\mathbf{x}) = \prod_{n=1}^N p(y_n|\mathbf{x})$$

$$p(y_n|\mathbf{x}) = p(y_n|\mathbf{x}^{[n]})$$

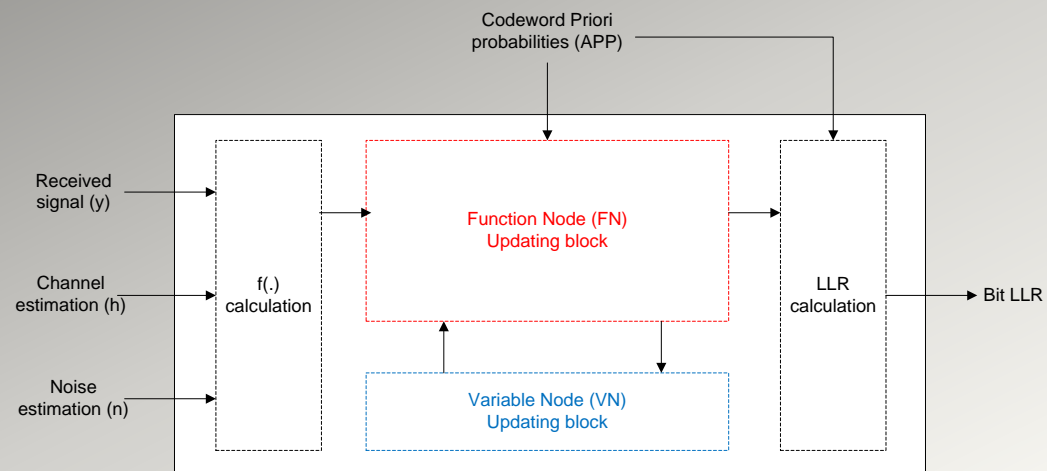
$$\hat{x}_k = \arg \max_{a \in \mathcal{X}} \sum_{\substack{\mathbf{x} \in \mathcal{X}^K \\ x_k = a}} P(\mathbf{x}) \prod_{n \in \zeta_k} p(y_n|\mathbf{x}^{[n]})$$

# How to Do SCMA Decoding with MPA Algorithm

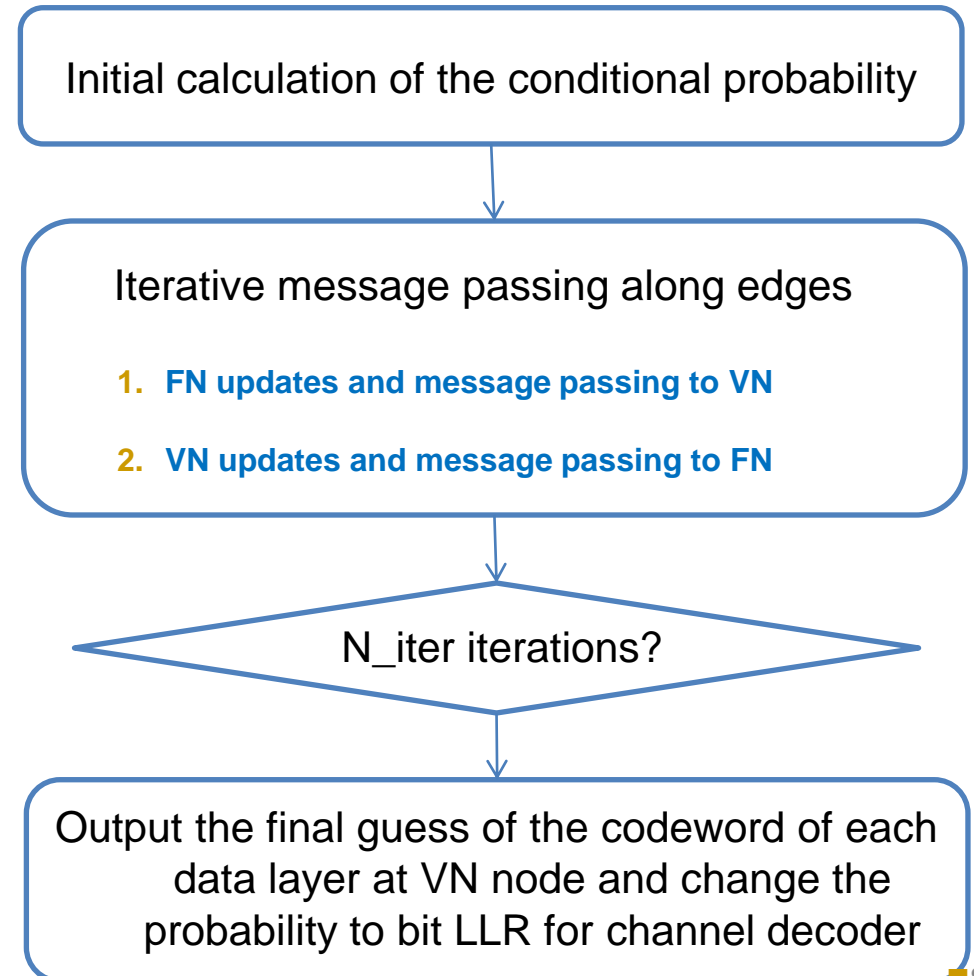
## Selection of SCMA Decoder

- The optimal multi-user detection can be done by using the maximum joint a posteriori probability (MAP) detection with excessive search – **non-practical complexity**
- With the low density spreading structure employed in SCMA, we can derive near ML performance multi-user detection with **message passing algorithm (MPA)** – **affordable complexity**

## Diagram for Message Passing Algorithm

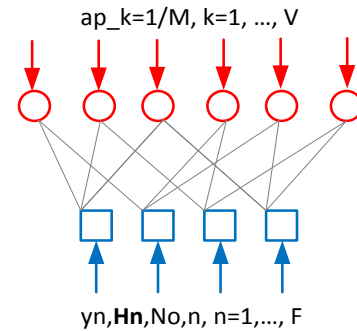


## MPA Decoder (Performed for each SCMA block)



# How to Do SCMA Decoding with MPA Algorithm

Parameters	Description of the parameters
$y_n, n=1, \dots, F$	Received signal as input to the MPA decoder on resource n
$m_k, k=1, \dots, V$	Codeword selected by layer k, $m_k = 1, \dots, M$
$N_{0,n}, n=1, \dots, F$	Noise power estimation on physical resource n
$C_{k,n}(m_k)$	The constellation symbol of VN node k on physical resource n when using codeword $m_k$
$H_n = \{h_{n,k}\}$	Channel gain of user k on physical resource n
$A_{p,k}, k=1, \dots, V$	A prior probability of codeword k, assuming equal probability $1/M$
$LLR_{k,b}$	logarithm of the likelihood ratio of layer k bit b
$N_{iter}$	Number of iterations in the MPA



## Step 1: Initial calculation of the conditional probability

- For each function node FN, calculate the  $f_n()$  function, which is the set of all possible residual signals given the known or estimated channel  $h_{n,k}$  and the assumed transmitted codeword  $C_{k,n}(m_k)$
- When  $d_f = 3$ , as in the example, for each FN node n, there are  $M^*M^*M$  combinations of transmitted signals, so there are in total  $F^*M^*M^*M$  values to store for  $f()$  function calculation

$$f_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n) = \frac{-1}{N_{0,n}} \left\| y_n - (h_{n,1} C_{1,n}(m_1) + h_{n,2} C_{2,n}(m_2) + h_{n,3} C_{3,n}(m_3)) \right\|^2$$

$$m_1 = 1, \dots, M \quad m_2 = 1, \dots, M \quad m_3 = 1, \dots, M \quad n = 1, \dots, F$$

- $\Phi_n()$  function is actually the conditional probability for given codeword combination, for Gaussian noise case, it is the exponential operation over  $f_n$  function, so the storage needed is the same

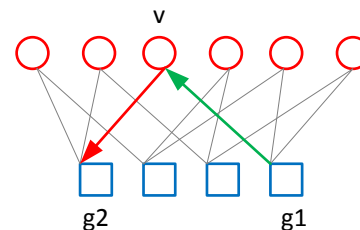
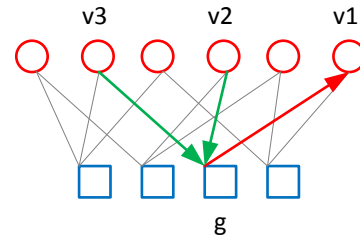
$$P(y_n | x_1, x_2, x_3) \text{ ----- } \Phi_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n) = \exp(f_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n))$$

- To prepare for the iterations, we assign the a prior probability for each codeword, which is assumed to be equal

$$P(x_1), P(x_2), P(x_3) \text{ ----- } I_{v_1 \rightarrow g}^{init}(m_1) = I_{v_2 \rightarrow g}^{init}(m_2) = I_{v_3 \rightarrow g}^{init}(m_3) = \frac{1}{M}$$

# How to Do SCMA Decoding with MPA Algorithm

Parameters	Description of the parameters
$y_n, n=1, \dots, F$	Received signal as input to the MPA decoder on resource n
$m_k, k=1, \dots, V$	Codeword selected by layer k, $m_k = 1, \dots, M$
$N_{0,n}, n=1, \dots, F$	Noise power estimation on physical resource n
$C_{k,n}(m_k)$	The constellation symbol of VN node k on physical resource n when using codeword $m_k$
$H_n = \{h_{n,k}\}$	Channel gain of user k on physical resource n
$A_{p,k}, k=1, \dots, V$	A prior probability of codeword k, assuming equal probability $1/M$
$LLR_{k,b}$	logarithm of the likelihood ratio of layer k bit b
$N_{iter}$	Number of iterations in the MPA



## Step 2: Iterative message passing along edges

### [FN update]: message passing from FN to its neighboring VNs

- FN node g passes updates obtained from extrinsic information to its neighboring VN nodes (g to v1, information from v2 and v3 are extrinsic)
- The message passed to v1 contains the guess of what signal at g may be given all possibilities of v1

$$I_{g \rightarrow v_1}(m_1) = \sum_{m_2=1}^M \sum_{m_3=1}^M \Phi_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n) (I_{v_2 \rightarrow g}(m_2) I_{v_3 \rightarrow g}(m_3)) \quad m_1 = 1, \dots, M$$

$$I_{g \rightarrow v_2}(m_2) = \sum_{m_1=1}^M \sum_{m_3=1}^M \Phi_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n) (I_{v_1 \rightarrow g}(m_1) I_{v_3 \rightarrow g}(m_3)) \quad m_2 = 1, \dots, M$$

$$I_{g \rightarrow v_3}(m_3) = \sum_{m_1=1}^M \sum_{m_2=1}^M \Phi_n(y_n, m_1, m_2, m_3, N_{0,n}, H_n) (I_{v_1 \rightarrow g}(m_1) I_{v_2 \rightarrow g}(m_2)) \quad m_3 = 1, \dots, M$$

### [VN update]: message passing from VN to its neighboring FNs

- VN node v passes updates obtained from extrinsic information to its neighboring FN nodes (v to g1, information from g2 is extrinsic)
- In the  $d_v=2$  case, it is actually a “guess” swap at VN node

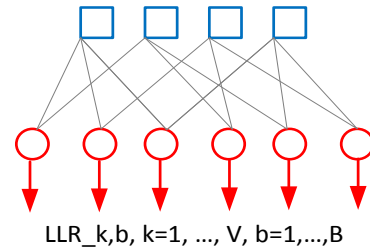
$$I_{v \rightarrow g_1}(m) = \text{normalize} (ap_v(m) I_{g_2 \rightarrow v}(m)) \quad m = 1, \dots, M$$

$$I_{v \rightarrow g_2}(m) = \text{normalize} (ap_v(m) I_{g_1 \rightarrow v}(m)) \quad m = 1, \dots, M$$



# How to Do SCMA Decoding with MPA Algorithm

Parameters	Description of the parameters
$y_n, n=1, \dots, F$	Received signal as input to the MPA decoder on resource n
$m_k, k=1, \dots, V$	Codeword selected by layer k, $m_k = 1, \dots, M$
$No_n, n=1, \dots, F$	Noise power estimation on physical resource n
$C_{k,n}(m_k)$	The constellation symbol of VN node k on physical resource n when using codeword $m_k$
$H_n = \{h_{n,k}\}$	Channel gain of user k on physical resource n
$Ap_k, k=1, \dots, V$	A prior probability of codeword k, assuming equal probability $1/M$
$LLR_{k,b}$	logarithm of the likelihood ratio of layer k bit b
$N_{iter}$	Number of iterations in the MPA



## Step 3: LLR output at variable node after $N_{iter}$ iterations

- After  $N_{iter}$  iterations, we shall output the guess at each VN node (for each data layer) as the detection results
- The guess at VN node  $v$  for codeword  $m$  is a chain product of all guesses from all its neighboring FN nodes and the a prior probability

$$Q_v(m) = ap_v(m) I_{g_1 \rightarrow v}(m) I_{g_2 \rightarrow v}(m) \quad m = 1, \dots, M$$

- After getting the probability guess of codeword at each layer, we then need to calculate the Log-Likelihood-Rate (LLR) for each coded bit, so that they can serve as the input for the turbo decoder (or any other channel decoder) directly after MPA

$$LLR_x = \log \left( \frac{P(b_x = 0)}{P(b_x = 1)} \right)$$

$$LLR_x = \log \left( \frac{\sum_{m: b_{m,x}=0} Q_v(m)}{\sum_{m: b_{m,x}=1} Q_v(m)} \right) = \log \left( \sum_{m: b_{m,x}=0} Q_v(m) \right) - \log \left( \sum_{m: b_{m,x}=1} Q_v(m) \right)$$

# Hints on Low Complexity MPA Receiver Design

- **Short-comings of the current MPA algorithm**

1. Though much lower complexity compared with the optimal MAP algorithm (thanks to the sparse structure of the SCMA codebook), it is still of high complexity for hardware
2. The  $\exp(\cdot)$  operations causes very large dynamic ranges and very high storage burden if using lookup table, which is not good news for hardware implementation

- **Hint 1: Change to LOG domain using Jacobi's logarithm**

1. After changing to Log domain,  $\exp(\cdot)$  operation disappears : **MPA -> MAX-Log MPA**

$$\log\left(\sum_{i=1}^N \exp(f_i)\right) \approx \max\{f_1, f_2, \dots, f_N\}$$

- **Hint 2: Optimize the calculations during iterations**

1. Try to optimize the order of iterations
2. Try to use as much as possible the common results in the calculation

To have a gut feeling how SCMA will be implemented in the 5G wireless systems

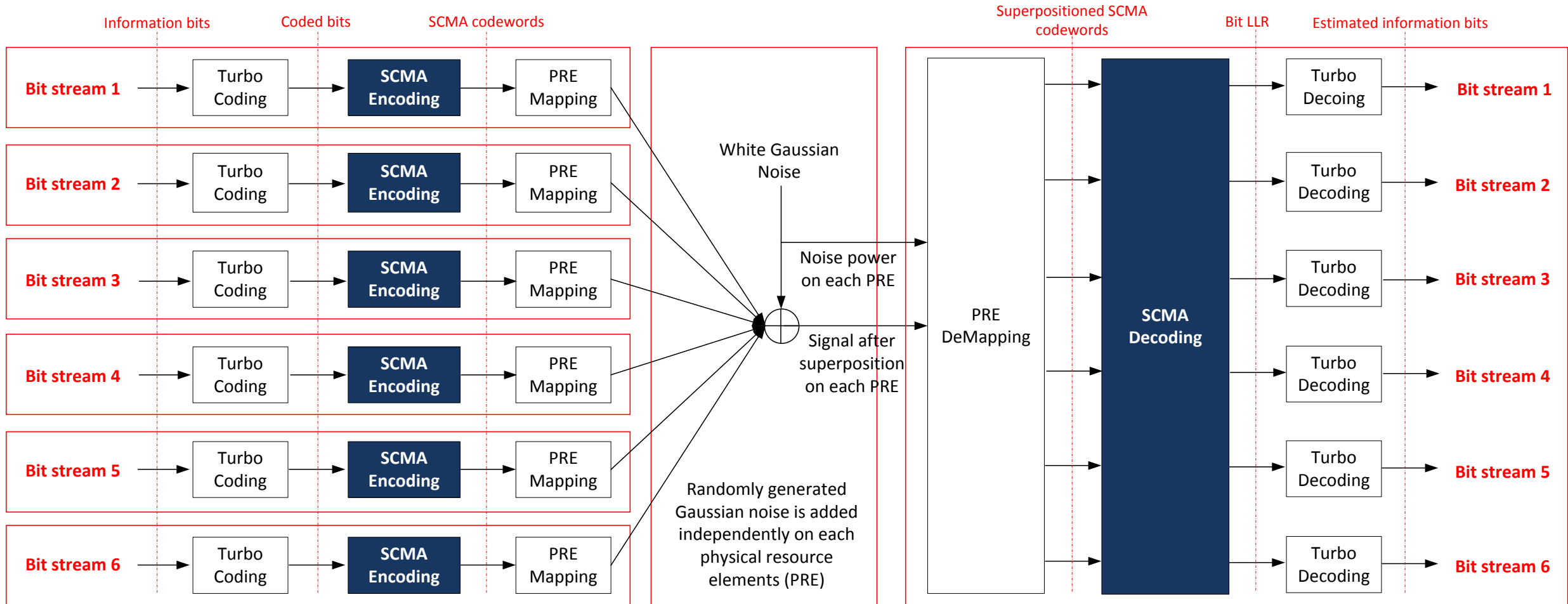
**WHAT WILL YOU IMPLEMENT?**

# Simplified Uplink SCMA System to be Implemented

Simplified Transmitter Side Representing 6 Users, each taking one layer of the SCMA codebook, namely CB<sub>i</sub> for user i

Simplified Communication Channel

Simplified Receiver Side Representing Base Station, with joint multi-user detection using MPA algorithm



\* Turbo coding and decoding can be replaced by other forward error correction (FEC) channel coding/decoding modules.

# System Configuration Parameters for Implementation

Parameter Categories	Related Variables	Typical value	Description
SCMA Codebook	V	6	6 variable nodes (VN), number of data layers
	F	4	4 function nodes (FN), number of physical resources
	d <sub>f</sub>	3	Each FN is connected to 3 VNs
	d <sub>v</sub>	2	Each VN is connected to 2 FNs
	M	4	Number of codeword in each codebook
	CB <sub>i</sub>	F-by-M matrix	Codebook for one SCMA data layer, given
Turbo Coding	R	1/2	Coding rate, defined as the ratio of information bits over coded bits
SCMA decoding	N <sub>iter</sub>	3 ~ 15	Number of iterations in MPA
	H <sub>n</sub>	{1}	Channel gain, in the white Gaussian noise only case, H <sub>n</sub> ={h <sub>n,k</sub> }= {1}
	APP <sub>i</sub>	1/M	A prior probability of codeword i, assuming equal probability 1/M
System Scale	B	125 bytes = 1000 bits	Total number of information bits, randomly generated
	N	B / R = 2000 bits	Total number of coded bits after Turbo coding
	L	L = F * N / log <sub>2</sub> (M) = 4000	Total number of physical resource units

# How We Judge and Compete the Results

- **Phase I with detailed design document and simulations**

**To deliver:**

1. Detailed design document for FPGA implementation
2. Matlab/C simulation code for the link and the BER v.s. Eb/No curve

**To check:**

1. Correct understanding of how SCMA system shall be implemented, including the SCMA encoder and the SCMA decoder
2. Low complexity design of SCMA decoder based on the hint given in the material, i.e., MAX-Log MPA

- **Phase II with complete FPGA implementation and test**

**To deliver:**

1. Complete FPGA implementation
2. BER v.s. Eb/No curves tested from FPGA implementation, should be align with simulation (1dB difference at most)

**To check:**

1. Bit streams can be decoded with the average bit-error-rate (BER) less than 0.001 (namely at most 1 bit error in the total 1000 bits)
2. FPGA resources used should be minimized through the design of low complexity SCMA decoder design and efficient way of code implementation

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